
*Density Elimination and Standard Completeness for
extensions of UL and MTL*

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Standard Completeness

Completeness of a logic with respect
to algebras whose lattice reduct is the real interval $[0, 1]$.

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Completeness of a logic with respect to algebras whose lattice reduct is the real interval $[0, 1]$.

- Intended semantics for *Fuzzy logic* (Hajek 1998)

Examples: standard complete logics

- HL : Logic of Continuous t-norms
- MTL: Logic of Left-continuous t-norms
- UL : Logic of Left-continuous uninorms

Our results

We prove standard completeness for

- Classes of axiomatic extensions of UL
- Classes of axiomatic extensions of MTL

The way to Standard Completeness

Given a logic L

1. General algebraic completeness, i.e. completeness w.r.t. L-algebras

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2. Completeness w.r.t. L -chains (linearly ordered L -algebras).

The way to Standard Completeness

Given a logic L

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2. Completeness w.r.t. L -chains.
 - $UL \iff UL$ -chains
 - $MTL \iff MTL$ -chains

The way to Standard Completeness

Given a logic L

1. General algebraic completeness, i.e. completeness w.r.t. L -algebras.

2. Completeness w.r.t. L -chains.

- $UL + \alpha \iff UL\text{-chains satisfying } 1 \leq \alpha$
- $MTL + \alpha \iff MTL\text{-chains satisfying } 1 \leq \alpha$

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3. Completeness w.r.t countable dense L -chains (rational completeness).

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The way to Standard Completeness

Given a logic L

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2. Completeness w.r.t. L -chains.
3. **Completeness w.r.t countable dense L -chains (rational completeness).**
 - (Metcalfe, Montagna JSL 2007)
Add the density rule to L

$$\frac{(\alpha \rightarrow p) \vee (p \rightarrow \beta) \vee \gamma}{(\alpha \rightarrow \beta) \vee \gamma} \text{ (density)}$$

$L + (\text{density})$ is rational complete.

Density Elimination and Rational completeness

$L + (density)$ is rational complete. We prove that L is rational complete as follows

- Find a suitable hypersequent calculus HL for L
- Show that the density rule is eliminable in HL

The way to Standard Completeness

Given a logic L

1. General algebraic completeness, i.e. completeness w.r.t. L -algebras.
2. Completeness w.r.t. L -chains.
3. Completeness of $L = L + (density)$ w.r.t countable dense L -chains (rational completeness).
4. Standard Completeness (via Dedekind-MacNeille completion)

Results on uninorm logic UL

$$UL = FL_e + ((\alpha \rightarrow \beta) \wedge t) \vee ((\beta \rightarrow \alpha) \wedge t)$$

We will show standard completeness for axiomatic extensions of UL with any *knotted axiom*, i.e.:

- $UL + (\alpha \rightarrow \alpha \cdot \alpha)$
- $UL + (\alpha \cdot \alpha \rightarrow \alpha)$
- $UL + (\alpha^k \rightarrow \alpha^n)$ (for $n, k > 1$)

Our basic calculus FL_e : sequent calculus

$$\frac{}{\alpha \Rightarrow \alpha} \text{ (init)}$$

$$\frac{}{\Rightarrow t} \text{ (tr)}$$

$$\frac{}{f \Rightarrow} \text{ (fl)}$$

$$\frac{}{\Gamma \Rightarrow \top} \text{ (}\top\text{)}$$

$$\frac{}{\Gamma, \perp \Rightarrow \Delta} \text{ (}\perp\text{)}$$

$$\frac{\Gamma \Rightarrow \Pi}{t, \Gamma \Rightarrow \Pi} \text{ (tl)}$$

$$\frac{\Gamma \Rightarrow}{\Gamma \Rightarrow f} \text{ (fr)}$$

$$\frac{\Gamma \Rightarrow \alpha \quad \alpha, \Delta \Rightarrow \Pi}{\Gamma, \Delta \Rightarrow \Pi} \text{ (cut)}$$

$$\frac{\Gamma \Rightarrow \alpha \quad \Gamma \Rightarrow \beta}{\Gamma \Rightarrow \alpha \wedge \beta} \text{ (}\wedge r\text{)}$$

$$\frac{\alpha_i, \Gamma \Rightarrow \Pi}{\alpha_1 \wedge \alpha_2, \Gamma \Rightarrow \Pi} \text{ (}\wedge l\text{)}$$

$$\frac{\Gamma \Rightarrow \alpha_i}{\Gamma \Rightarrow \alpha_1 \vee \alpha_2} \text{ (}\vee r\text{)}$$

$$\frac{\alpha, \Gamma \Rightarrow \Pi \quad \beta, \Gamma \Rightarrow \Pi}{\alpha \vee \beta, \Gamma \Rightarrow \Pi} \text{ (}\vee l\text{)}$$

$$\frac{\Gamma \Rightarrow \alpha \quad \beta, \Delta \Rightarrow \Pi}{\Gamma, \alpha \rightarrow \beta, \Delta \Rightarrow \Pi} \text{ (}\rightarrow l\text{)}$$

$$\frac{\alpha, \Gamma \Rightarrow \beta}{\Gamma \Rightarrow \alpha \rightarrow \beta} \text{ (}\rightarrow r\text{)}$$

$$\frac{\Gamma \Rightarrow \alpha \quad \Delta \Rightarrow \beta}{\Gamma, \Delta \Rightarrow \alpha \cdot \beta} \text{ (}\cdot r\text{)}$$

$$\frac{\alpha, \beta, \Gamma \Rightarrow \Pi}{\alpha \cdot \beta, \Gamma \Rightarrow \Pi} \text{ (}\cdot l\text{)}$$

Hypersequent calculus HUL for UL

(Avron '89): Hypersequent

$$\Gamma_1 \Rightarrow \Pi_1 \mid \dots \mid \Gamma_n \Rightarrow \Pi_n$$

where for all $i = 1, \dots, n$, $\Gamma_i \Rightarrow \Pi_i$ is an ordinary sequent
| is intended to denote a meta-level disjunction.

Hypersequent calculus HUL for UL

Embed sequent rules for FL_e into hypersequents

$$\begin{array}{ccc}
 \frac{}{G \Rightarrow t} \text{ (tr)} & \frac{}{G|\alpha \Rightarrow \alpha} \text{ (init)} & \frac{}{G|f \Rightarrow} \text{ (fl)} \\
 \\
 \frac{}{G|\Gamma \Rightarrow \top} \text{ (T)} & \frac{}{G|\Gamma, \perp \Rightarrow \Delta} \text{ (\perp)} & \\
 \\
 \frac{G|\Gamma \Rightarrow \alpha \quad G|\alpha, \Delta \Rightarrow \Pi}{G|\Gamma, \Delta \Rightarrow \Pi} \text{ (cut)} & \frac{G|\Gamma \Rightarrow \Pi}{G|t, \Gamma \Rightarrow \Pi} \text{ (tl)} & \frac{G|\Gamma \Rightarrow}{G|\Gamma \Rightarrow f} \text{ (fr)} \\
 \\
 \frac{G|\Gamma \Rightarrow \alpha \quad G|\Gamma \Rightarrow \beta}{G|\Gamma \Rightarrow \alpha \wedge \beta} \text{ (\wedge r)} & \frac{G|\alpha_i, \Gamma \Rightarrow \Pi}{G|\alpha_1 \wedge \alpha_2, \Gamma \Rightarrow \Pi} \text{ (\wedge l)} & \frac{G|\Gamma \Rightarrow \alpha_i}{G|\Gamma \Rightarrow \alpha_1 \vee \alpha_2} \text{ (\vee r)} \\
 \\
 \frac{G|\alpha, \Gamma \Rightarrow \Pi \quad G|\beta, \Gamma \Rightarrow \Pi}{G|\alpha \vee \beta, \Gamma \Rightarrow \Pi} \text{ (\vee l)} & \frac{G|\Gamma \Rightarrow \alpha \quad G|\beta, \Delta \Rightarrow \Pi}{G|\Gamma, \alpha \rightarrow \beta, \Delta \Rightarrow \Pi} \text{ (\rightarrow l)} & \frac{G|\alpha, \Gamma \Rightarrow \beta}{G|\Gamma \Rightarrow \alpha \rightarrow \beta} \text{ (\rightarrow r)} \\
 \\
 \frac{G|\Gamma \Rightarrow \alpha \quad G|\Delta \Rightarrow \beta}{G|\Gamma, \Delta \Rightarrow \alpha \cdot \beta} \text{ (\cdot r)} & \frac{G|\alpha, \beta, \Gamma \Rightarrow \Pi}{G|\alpha \cdot \beta, \Gamma \Rightarrow \Pi} \text{ (\cdot l)} &
 \end{array}$$

Hypersequent calculus HUL for UL

We add:

- Suitable rules to manipulate the additional layer of structure.

$$\frac{G}{G | \Gamma \Rightarrow \alpha} \text{ (ew)}$$

$$\frac{G | \Gamma \Rightarrow \alpha | \Gamma \Rightarrow \alpha}{G | \Gamma \Rightarrow \alpha} \text{ (ec)}$$

Hypersequent calculus HUL for UL

We add:

- Suitable rules to manipulate the additional layer of structure.

$$\frac{G}{G | \Gamma \Rightarrow \alpha} \text{ (ew)}$$

$$\frac{G | \Gamma \Rightarrow \alpha \mid \Gamma \Rightarrow \alpha}{G | \Gamma \Rightarrow \alpha} \text{ (ec)}$$

- A hypersequent structural rule corresponding to prelinearity :

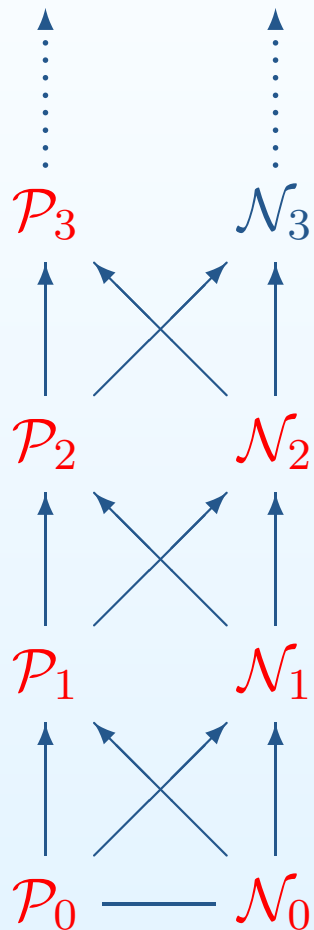
$$((\alpha \rightarrow \beta) \wedge t) \vee ((\beta \rightarrow \alpha) \wedge t)$$

$$\frac{G | \Gamma_1, \Delta_1 \Rightarrow \Pi_1 \quad G | \Gamma_2, \Delta_2 \Rightarrow \Pi_2}{G | \Gamma_1, \Gamma_2 \Rightarrow \Pi_1 \mid \Delta_1, \Delta_2 \Rightarrow \Pi_2} \text{ (com)}$$

Hypersequent calculi for extensions of UL ?

(Ciabattoni, Galatos, Terui 2008).

Sets $\mathcal{P}_n, \mathcal{N}_n$ of formulas defined by:



$\mathcal{P}_0, \mathcal{N}_0 :=$ Atomic formulas

$\mathcal{P}_{n+1} := \mathcal{N}_n \mid \mathcal{P}_{n+1} \cdot \mathcal{P}_{n+1} \mid \mathcal{P}_{n+1} \vee \mathcal{P}_{n+1} \mid 1 \mid \perp$

$\mathcal{N}_{n+1} := \mathcal{P}_n \mid \mathcal{P}_{n+1} \rightarrow \mathcal{N}_{n+1} \mid \mathcal{N}_{n+1} \wedge \mathcal{N}_{n+1} \mid 0 \mid \top$

Examples:

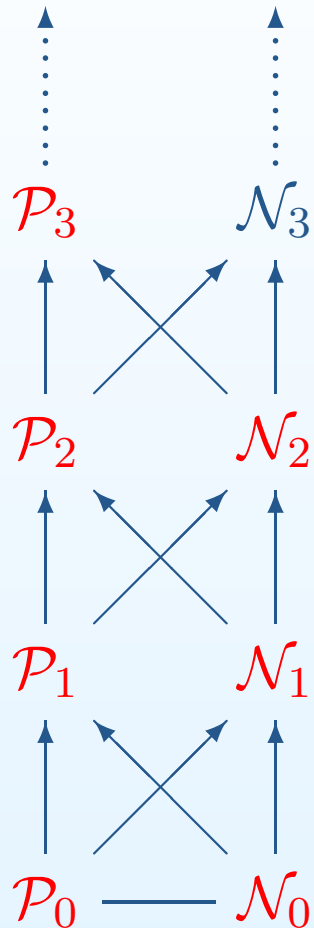
- To the class \mathcal{N}_2 belong :

$$\alpha \rightarrow \alpha \cdot \alpha \quad \alpha \cdot \alpha \rightarrow \alpha \quad \alpha^k \rightarrow \alpha^n$$

- To the class \mathcal{P}_3 belong :

$$\neg \alpha \vee \neg \neg \alpha \quad \neg(\alpha \cdot \beta) \vee ((\alpha \wedge \beta) \rightarrow (\alpha \cdot \beta))$$

Hypersequent calculi for extensions of UL ?



Algorithm to convert axioms into “good” rules, preserving cut-elimination.

- Axioms in $\mathcal{N}_2 \Rightarrow$ Sequent structural rules
- Axioms in (subclass of) $\mathcal{P}_3 \Rightarrow$ Hypersequent structural rules

Correspondence axioms - rules

Class	Axiom	Rule
\mathcal{N}_2	$(\alpha \rightarrow t) \wedge (f \rightarrow \alpha)$	$\frac{G \Pi \Rightarrow \Psi}{G \Pi, \alpha \Rightarrow \Psi} \quad (wl) \quad \frac{G \Pi \Rightarrow \Psi}{G \Pi \Rightarrow \alpha} \quad (wr)$
	$\alpha \rightarrow \alpha \cdot \alpha$	$\frac{G_1 \Pi, \Gamma, \Gamma \Rightarrow \Psi}{G_1 \Pi, \Gamma \Rightarrow \Psi} \quad (c)$
	$\alpha \cdot \alpha \rightarrow \alpha$	$\frac{G_1 \Pi, \Gamma_1 \Rightarrow \Psi \quad G_1 \Pi, \Gamma_2 \Rightarrow \Psi}{G_1 \Pi, \Gamma_1, \Gamma_2, \Rightarrow \Psi} \quad (mgl)$
	$\alpha^k \rightarrow \alpha^n$	$\frac{G \Pi, \Gamma_1^n \Rightarrow \Psi \quad \dots \quad G_1 \Pi, \Gamma_k^n \Rightarrow \Psi}{G_1 \Pi, \Gamma_1, \dots, \Gamma_k \Rightarrow \Psi} \quad (knot_k^n)$

Density elimination

- Density rule in hypersequent calculus :

$$\frac{G \mid \Gamma \Rightarrow p \mid p \Rightarrow \Delta}{G \mid \Gamma \Rightarrow \Delta} \text{ (density)}$$

where p does not occur in the conclusion (*eigenvariable*).

Density elimination

- Density rule in hypersequent calculus :

$$\frac{G \mid \Gamma \Rightarrow p \mid p \Rightarrow \Delta}{G \mid \Gamma \Rightarrow \Delta} \text{ (density)}$$

where p does not occur in the conclusion (*eigenvariable*).

- Similar to cut elimination

$$\frac{G \mid \Gamma \Rightarrow A \quad G \mid A, \Sigma \Rightarrow \Delta}{G \mid \Gamma, \Sigma \Rightarrow \Delta} \text{ (cut)}$$

Proof by induction on the length of derivations

Density elimination

(Ciabattoni, Metcalfe TCS 2008)

Given a density-free derivation, ending in

$$\frac{\begin{array}{c} \vdots d \\ G \mid \Gamma \Rightarrow p \mid p \Rightarrow \Delta \end{array}}{G \mid \Gamma \Rightarrow \Delta} \text{ (density)}$$

Density elimination

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$$\frac{\begin{array}{c} \vdots d \\ G \mid \Gamma \Rightarrow p \mid p \Rightarrow \Delta \end{array}}{G \mid \Gamma \Rightarrow \Delta} \text{ (density)}$$

- **Asymmetric substitution:** p is replaced
 - With Δ when occurring on the right
 - With Γ when occurring on the left

Density elimination

$$\frac{\begin{array}{c} \vdots d^* \\ G \mid \Gamma \Rightarrow \Delta \mid \Gamma \Rightarrow \Delta \end{array}}{G \mid \Gamma \Rightarrow \Delta} \text{ (EC)}$$

- **Asymmetric substitution:** p is replaced
 - With Δ when occurring on the right
 - With Γ when occurring on the left

A problem

$$\frac{\begin{array}{c} p \Rightarrow p \\ \vdots \\ d \\ \vdots \end{array} \quad G \mid \Gamma \Rightarrow p \mid p \Rightarrow \Delta}{G \mid \Gamma \Rightarrow \Delta} \text{ (EC)}$$

- $p \Rightarrow p$ axiom

A problem

$$\frac{\begin{array}{c} \Gamma \Rightarrow \Delta \\ \vdots \\ d \\ \vdots \\ G \mid \Gamma \Rightarrow \Delta \mid \Gamma \Rightarrow \Delta \end{array}}{G \mid \Gamma \Rightarrow \Delta} \text{ (EC)}$$

- $p \Rightarrow p$ axiom
- $\Gamma \Rightarrow \Delta$ not an axiom

A possible solution

$$\frac{\begin{array}{c} p \Rightarrow p \\ \vdots \\ d \\ \vdots \end{array}}{G \mid \Gamma \Rightarrow p \mid p \Rightarrow \Delta} \text{(EC)} \\ G \mid \Gamma \Rightarrow \Delta$$

A possible solution

(Ciabattoni, Metcalfe 2008)

$$\frac{\begin{array}{c} \Rightarrow t \\ \vdots \\ d^* \\ \vdots \\ G \mid \Gamma \Rightarrow \Delta \mid \Gamma \Rightarrow \Delta \end{array}}{G \mid \Gamma \Rightarrow \Delta} \text{ (EC)}$$

- We substitute:
 - $p \Rightarrow p$ (axiom) with $\Rightarrow t$ (axiom).
 - p with Δ when occurring on the right.
 - p with Γ when occurring on the left.

Works for *HUL*. What about extensions?

Contraction and mingle

Consider UL extended with

-

$$\alpha \rightarrow \alpha \cdot \alpha$$

-

$$\alpha \cdot \alpha \rightarrow \alpha$$

Contraction and mingle

Hypersequent calculus: *HUL* plus

-

$$\frac{G_1 | \Pi, \Gamma, \Gamma \Rightarrow \Psi}{G_1 | \Pi, \Gamma \Rightarrow \Psi} \quad (c)$$

-

$$\frac{G_1 | \Pi, \Gamma_1 \Rightarrow \Psi \quad G_1 | \Pi, \Gamma_2 \Rightarrow \Psi}{G_1 | \Pi, \Gamma_1, \Gamma_2 \Rightarrow \Psi} \quad (mgl)$$

Recall

(Ciabattoni, Metcalfe 2008)

$$\frac{\begin{array}{c} \Rightarrow t \\ \vdots \\ d^* \\ \vdots \\ G \mid \Gamma \Rightarrow \Delta \mid \Gamma \Rightarrow \Delta \end{array}}{G \mid \Gamma \Rightarrow \Delta} \text{ (EC)}$$

- We substitute:
 - $p \Rightarrow p$ (axiom) with $\Rightarrow t$ (axiom).
 - p with Δ when occurring on the right.
 - p with Γ when occurring on the left.

Problematic case

Consider

$$\frac{\Pi, p, p \Rightarrow p}{\Pi, p \Rightarrow p} \quad (c)$$

Can we get:

$$\frac{\Pi, \Gamma \Rightarrow t}{\Pi \Rightarrow t} \quad ?$$

Proof by cases

Recall: the hypersequent symbol '|' is interpreted as disjunction.

- For any hypersequent calculus HL the following meta-rule holds:

$$\frac{G_1 \vdash_{HL} H \quad G_2 \vdash_{HL} H}{G_1 | G_2 \vdash_{HL} H}$$

Proof by cases and Density Elimination

Recall our derivation

$$\frac{\begin{array}{c} \vdots d \\ G \mid \Gamma \Rightarrow p \mid p \Rightarrow \Delta \end{array}}{G \mid \Gamma \Rightarrow \Delta} \text{ (density)}$$

Proof by cases and Density Elimination

In d we instantiate p with t , obtaining

$$G \mid \Gamma \Rightarrow \begin{array}{c} \vdots \\ d_t \\ \vdots \end{array} t \mid t \Rightarrow \Delta$$

Proof by cases and Density Elimination

We instantiate p with t , obtaining

$$G \mid \Gamma \Rightarrow \begin{array}{c} \vdots d_t \\ t \mid t \end{array} \Rightarrow \Delta$$

We find density free proofs of:

$$\begin{array}{c} G \mid \Gamma \Rightarrow t \\ \vdots d_1 \\ G \mid \Gamma \Rightarrow \Delta \end{array}$$

$$\begin{array}{c} G \mid t \Rightarrow \Delta \\ \vdots d_2 \\ G \mid \Gamma \Rightarrow \Delta \end{array}$$

Proof by cases and Density Elimination

We find density free proofs of:

$$\begin{array}{l} G|\Gamma \Rightarrow t \\ \quad \vdots d_1 \\ G|\Gamma \Rightarrow \Delta \end{array} \qquad \begin{array}{l} G|t \Rightarrow \Delta \\ \quad \vdots d_2 \\ G|\Gamma \Rightarrow \Delta \end{array}$$

Applying the proof by cases property, we get:

$$\begin{array}{l} G|\Gamma \Rightarrow t|t \Rightarrow \Delta \\ \quad \vdots \\ G|\Gamma \Rightarrow \Delta \end{array}$$

Proof by cases and Density Elimination

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Applying the proof by cases property, we get:

$$\begin{array}{l} \quad \quad \quad \vdots d_t \\ G|\Gamma \Rightarrow t|t \Rightarrow \Delta \\ \quad \quad \quad \vdots \\ G|\Gamma \Rightarrow \Delta \end{array}$$

Knotted rules

- Consider UL extended with

$$\alpha^k \rightarrow \alpha^n$$

for $k, n > 1$

Knotted rules

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$$\alpha^k \rightarrow \alpha^n$$

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- Hypersequent calculus: HUL plus

$$\frac{G|\Pi, \Gamma_1^n \Rightarrow \Psi \quad \dots \quad G_1|\Pi, \Gamma_k^n \Rightarrow \Psi}{G_1|\Pi, \Gamma_1, \dots, \Gamma_k \Rightarrow \Psi} \text{ (knot}_k^n\text{)}$$

Knotted rules: Problematic case

$$\frac{\Pi, p, p, p \Rightarrow p \quad \Pi, p, p, p \Rightarrow p}{\Pi, p, p \Rightarrow p} \text{ (knot}_2^3\text{)}$$

We would like to get:

$$\frac{\Pi, \Gamma, \Gamma \Rightarrow t \quad \Pi, \Gamma, \Gamma \Rightarrow t}{\Pi, \Gamma \Rightarrow t} ?$$

Knotted rules: Problematic case

$$\frac{\Pi, p, p, p \Rightarrow p \quad \Pi, p, p, p \Rightarrow p}{\Pi, p, p \Rightarrow p} \text{ (knot}_2^3\text{)}$$

Using some derivabilities in $HUL + (\text{knot}_2^3)$, we can show:

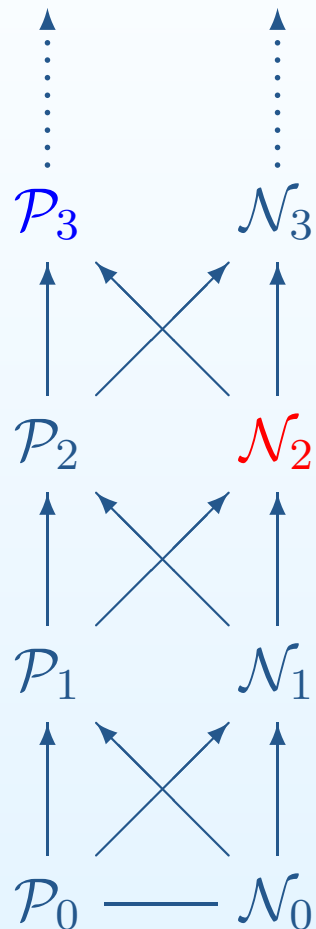
$$\begin{array}{c} \Pi, \Gamma, \Gamma \Rightarrow t \quad \Pi, \Gamma, \Gamma \Rightarrow t \\ \vdots \\ \Pi, \Gamma \Rightarrow t \end{array}$$

Standard completeness for extensions of UL

Let $L = UL + (\alpha^k \rightarrow \alpha^n)$

1. General algebraic completeness, i.e. completeness w.r.t. L -algebras.
2. Completeness w.r.t. L -chains.
3. $L = L + (density)$ is Rational complete .
4. L is Standard complete (via Dedekind-MacNeille completion).

Closure under DM-completions



(Ciabattoni, Terui, Galatos 2011) Axioms on FL \leftrightarrow equations over residuated lattices

- A subclass of equations in class \mathcal{N}_2 are preserved by Dedekind-MacNeille completion. All the axioms we considered are in this class.
- A subclass of equations in class \mathcal{P}_3 are preserved by Dedekind-MacNeille completion, when applied to subdirectly irreducible algebras

Standard completeness for extensions of MTL

- $MTL = UL + (f \rightarrow \alpha) \wedge (\alpha \rightarrow t)$
- Hypersequent calculus $HMTL = HUL + (wl) + (wr)$

$$\frac{G \mid \Pi \Rightarrow \Psi}{G \mid \Pi, \alpha \Rightarrow \Psi} (wl) \qquad \frac{G \mid \Pi \Rightarrow}{G \mid \Pi \Rightarrow \alpha} (wr)$$

Standard completeness for extensions of MTL

- Density Elimination holds for $HMTL$ extended with any structural sequent rule
 - Any axiomatic extension of MTL with axioms within \mathcal{N}_2 is standard complete (2008 Ciabattoni, Metcalfe).

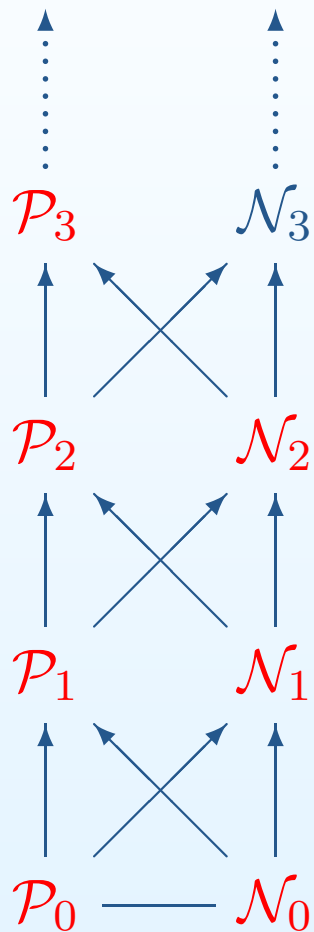
Standard completeness for extensions of *MTL*

- Density Elimination holds for *HMTL* extended with any structural sequent rule
 - Any axiomatic extension of *MTL* with axioms within \mathcal{N}_2 is standard complete (2008 Ciabattoni, Metcalfe).
- Density elimination holds for extensions of *HMTL* with structural hypersequent rules which do not “mix too much” the components (convergent rules).
 - Any axiomatic extension of *MTL* with axioms within a subclass of \mathcal{P}_3 is standard complete (2012 Baldi, Ciabattoni, Spendier)

Recall: Correspondence axioms-rules

(Ciabattoni, Galatos, Terui 2008).

Sets $\mathcal{P}_n, \mathcal{N}_n$ of formulas defined by:



$\mathcal{P}_0, \mathcal{N}_0 :=$ Atomic formulas

$\mathcal{P}_{n+1} := \mathcal{N}_n \mid \mathcal{P}_{n+1} \cdot \mathcal{P}_{n+1} \mid \mathcal{P}_{n+1} \vee \mathcal{P}_{n+1} \mid 1 \mid \perp$

$\mathcal{N}_{n+1} := \mathcal{P}_n \mid \mathcal{P}_{n+1} \rightarrow \mathcal{N}_{n+1} \mid \mathcal{N}_{n+1} \wedge \mathcal{N}_{n+1} \mid 0 \mid \top$

Examples:

- To the class \mathcal{N}_2 belong :

$$\alpha \rightarrow \alpha \cdot \alpha \quad \alpha \cdot \alpha \rightarrow \alpha$$

- To the class \mathcal{P}_3 belong :

$$\neg \alpha \vee \neg \neg \alpha \quad ((\alpha \rightarrow \beta) \wedge t) \vee ((\beta \rightarrow \alpha) \wedge t)$$

Recall: Correspondence axioms - rules

Class	Axiom	Rule
\mathcal{N}_2	$(\alpha \rightarrow t) \wedge (f \rightarrow \alpha)$ $\alpha \rightarrow \alpha \cdot \alpha$ $\alpha \cdot \alpha \rightarrow \alpha$ $\alpha^k \rightarrow \alpha^n$	$\frac{G \mid \Pi \Rightarrow \Psi}{G \mid \Pi, \alpha \Rightarrow \Psi} \text{ (wl)} \quad \frac{G \mid \Pi \Rightarrow}{G \mid \Pi \Rightarrow \alpha} \text{ (wr)}$ $\frac{G_1 \mid \Pi, \Gamma, \Gamma \Rightarrow \Psi}{G_1 \mid \Pi, \Gamma \Rightarrow \Psi} \text{ (c)}$ $\frac{G_1 \mid \Pi, \Gamma_1 \Rightarrow \Psi \quad G_1 \mid \Pi, \Gamma_2 \Rightarrow \Psi}{G_1 \mid \Pi, \Gamma_1, \Gamma_2, \Rightarrow \Psi} \text{ (mgl)}$ $\frac{G \mid \Pi, \Gamma_1^n \Rightarrow \Psi \quad \dots \quad G_1 \mid \Pi, \Gamma_k^n \Rightarrow \Psi}{G_1 \mid \Pi, \Gamma_1, \dots, \Gamma_k \Rightarrow \Psi} \text{ (knot}_k^n)$
\mathcal{P}_2	$\alpha \vee \neg \alpha$	$\frac{G \mid \Pi, \Gamma \Rightarrow \Psi}{G \mid \Gamma \Rightarrow \mid \Pi \Rightarrow \Psi} \text{ (em)}$
\mathcal{P}_3	$\neg \alpha \vee \neg \neg \alpha$	$\frac{G \mid \Gamma_1, \Gamma_2 \Rightarrow}{G \mid \Gamma_1 \Rightarrow \mid \Gamma_2 \Rightarrow} \text{ (lq)}$

Examples of convergent rules

- Axioms in \mathcal{P}_3 extending *MTL*

- (*wnm*) :

$$\neg(\alpha \cdot \beta) \vee ((\alpha \wedge \beta) \rightarrow (\alpha \cdot \beta))$$

- (*lq*) :

$$\neg\alpha \vee \neg\neg\alpha$$

- Corresponding convergent rules

$$\circ \frac{G \mid \Gamma_2, \Gamma_1, \Delta_1 \Rightarrow \Pi_1 \quad G \mid \Gamma_1, \Gamma_3, \Delta_1 \Rightarrow \Pi_1}{G \mid \Gamma_2, \Gamma_3 \Rightarrow \mid \Gamma_1, \Delta_1 \Rightarrow \Pi_1} \text{ (wnm)}$$

$$\circ \frac{G \mid \Gamma_1, \Gamma_2 \Rightarrow}{G \mid \Gamma_1 \Rightarrow \mid \Gamma_2 \Rightarrow} \text{ (lq)}$$

A non convergent rule

- Axiom in \mathcal{P}_3 extending *MTL*

- $\alpha \vee \neg\alpha$

- Corresponding rule

-

$$\frac{G|\Gamma, \Sigma \Rightarrow \Delta}{G|\Gamma \Rightarrow \mid \Sigma \Rightarrow \Delta} \text{ (em)}$$

Our results

- Standard completeness for extensions of UL :
 - $UL + \alpha^k \rightarrow \alpha^n$ (includes mingle and contraction axioms).
- Standard completeness for extensions of MTL :
 - Any axiomatic extension of MTL with axioms within a subclass of \mathcal{P}_3 is standard complete.

Open problem

- Conjecture: Let α be any formula in \mathcal{N}_2 . $UL + \alpha$ is standard complete.
- Let (r) be any sequent rule obtained by an axiom in \mathcal{N}_2 . Does $HUL + (r)$ admits Density Elimination?
 - Example, what about:

$$\frac{G \mid \Gamma_1 \Rightarrow \Gamma_2^2, \Pi \Rightarrow \Psi}{G \mid \Gamma_1, \Gamma_2, \Pi \Rightarrow \Psi}$$

Work in progress

- A general characterization of density elimination, hence standard completeness, for:
 - Extensions of MTL with axioms up to the class \mathcal{P}_3 in the substructural hierarchy.
 - Extensions of UL with axioms up to the class \mathcal{N}_2 in the substructural hierarchy.
 - Noncommutative variants of MTL and UL .
 - Logics with involutive negation. Long standing open problem: IUL

Appendix A: A class of structural rules

Let HL be HUL extended with any structural sequent rule

$$\frac{G_1 | \Pi_1, \Psi_1 \Rightarrow \Delta_1 \quad \dots \quad \Pi_1, \Psi_m \Rightarrow \Delta_1}{G_1 | \Pi_1, \Gamma_1, \dots, \Gamma_k \Rightarrow \Delta_1} \quad (r)$$

HL admits density elimination if (r) satisfies the following:

- Each Ψ_i is a multiset $\{\Gamma_{i_1}, \dots, \Gamma_{i_{n_i}}\}$ with $i_1 \dots i_{n_i}$ varying over $\{1, \dots, k\}$
- Either the minimum among the n_i is bigger than k or the maximum is smaller than k
- For any Γ_i there is at least one Ψ_j where Γ_i does not appear.
- For any Γ_i there is at least one Ψ_j where Γ_i appears more than once .

Appendix B: Convergent rules

Definition. Let (r) be a hypersequent structural rule with $G|S_i, i \in \{1, \dots, m\}$ premises, $C_1 | \dots | C_q$ conclusion.

- (0-pivot) $G|S_i$ is a *0-pivot* if there is an $s \in \{1, \dots, q\}$ such that $R(S_i) = R(C_s)$ and metavariables in $L(S_i)$ are contained in $L(C_s)$.
- (n-pivot) $G|S_j$ is an *n-pivot for $G|S_i$* with respect to $[\Delta_k / \Gamma_k]_{k \in \{1, \dots, n\}}$, with $\Gamma_k \in L(S_i)$ and $\Delta_k \in L(S_j)$, if the following conditions hold:
- $G|S_j$ is a *0-pivot*
 - $R(S_i) = R(S_j)$,
 - $L(S_j) = L(S_i[\Delta_k / \Gamma_k]_{k \in \{1, \dots, n\}})$,
 - If $n > 1$, $G|S_j$ is a $(n - 1)$ -*pivot* for n premises $G|S_{j_p}, p = 1, \dots, n$, with respect to $[\Delta_k / \Gamma_k]_{k \in \{1, \dots, n\} \setminus \{p\}}$.

Definition. A completed hypersequent rule (r) is *convergent* if for each premise $G|S_i$ one of the following conditions holds:

- $R(S_i) = \emptyset$,
- $G|S_i$ is a *0-pivot*
- there is a premise $G|S_j$ which is an *n-pivot for $G|S_i$* , with $n > 0$.