

# ICANN 19

## Counting with Analog Neurons

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## Motivations: The Computational Power of NNs

(discrete-time recurrent NNs with the saturated-linear activation function)

depends on the information contents of weight parameters:

1. **integer** weights: **finite automaton** (Minsky, 1967)
2. **rational** weights: **Turing machine** (Siegelmann, Sontag, 1995)  
polynomial time  $\equiv$  complexity class P  
  
polynomial time & increasing **Kolmogorov complexity** of real weights  $\equiv$   
a proper **hierarchy** of nonuniform complexity classes between P and P/poly  
(Balcázar, Gavalda, Siegelmann, 1997)
3. arbitrary **real** weights: **“super-Turing” computation** (Siegelmann, Sontag, 1994)  
polynomial time  $\equiv$  nonuniform complexity class P/poly  
exponential time  $\equiv$  any I/O mapping

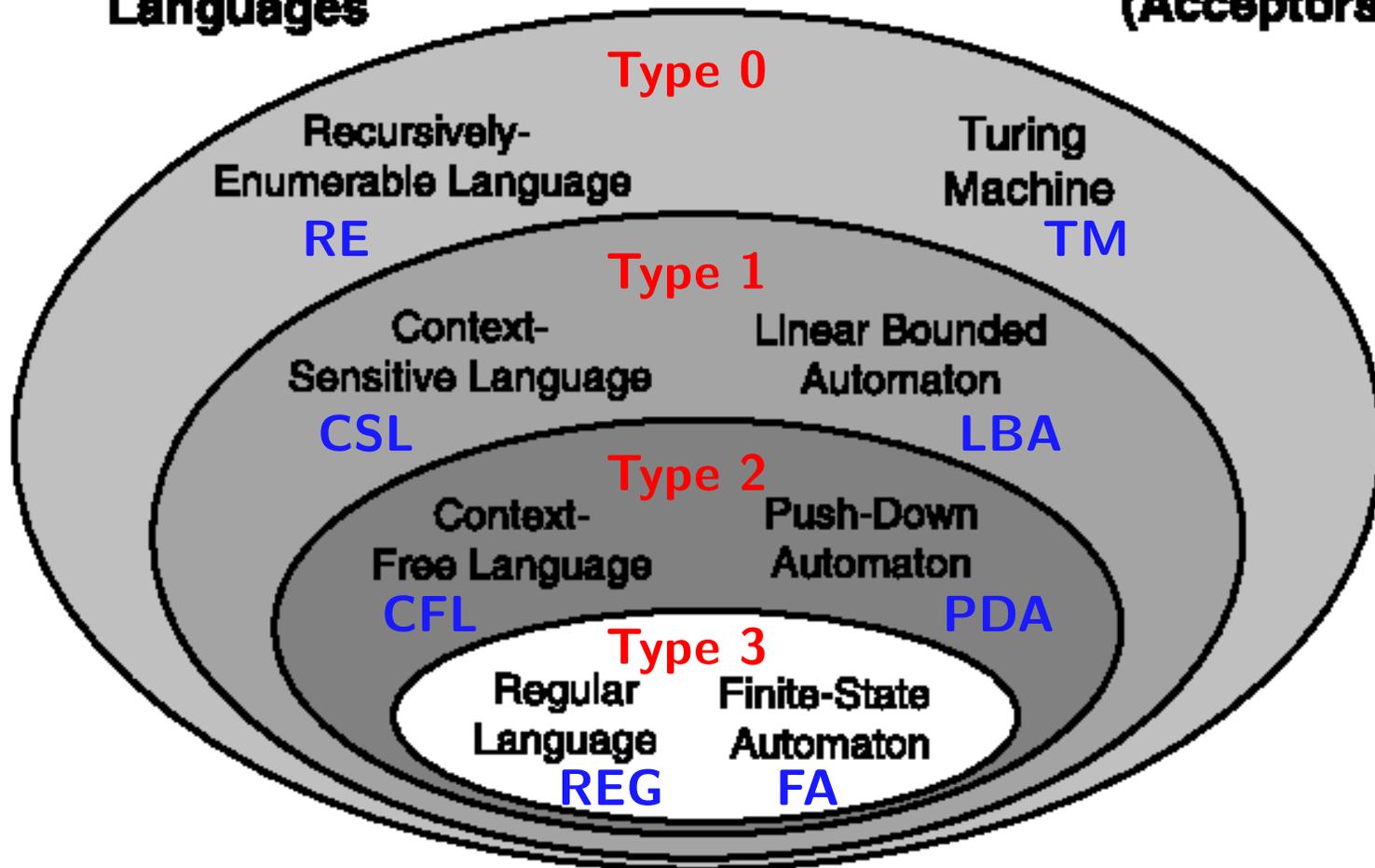
**filling the gap** between **integer** and **rational** weights w.r.t. **the Chomsky hierarchy**

regular (Type 3)  $\times$  recursively enumerable (Type 0) languages

# The Traditional Chomsky Hierarchy

**Grammars (Generators) &  
Languages**

**Automata  
(Acceptors)**



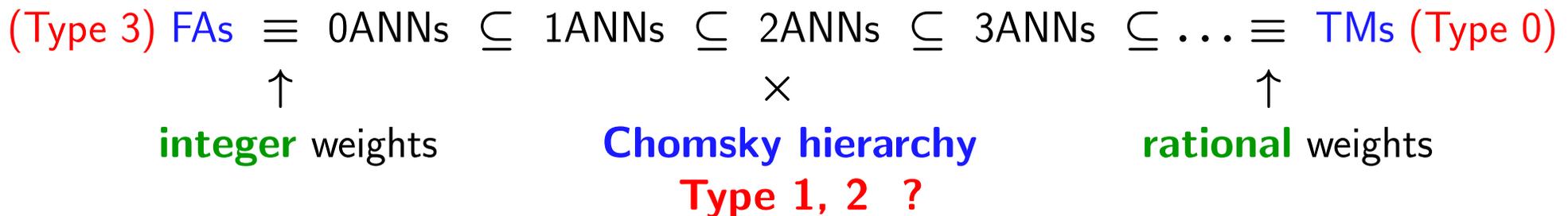
**The Formal Language Hierarchy**

# The Analog Neuron Hierarchy

$\alpha$ ANN = a binary-state NN with **integer** weights

+  $\alpha$  **extra analog-state neurons** with **rational** weights

the computational power of NNs increases with the number  $\alpha$  of extra analog neurons:



## Known Results:

- classifying **1ANNs** within the Chomsky hierarchy (Šíma, 2017):
  - upper bound: **1ANNs**  $\subset$  **LBAs**  $\equiv$  **CSLs (Type 1)**
  - lower bound: **1ANNs**  $\not\subset$  **PDA**s  $\equiv$  **CFLs (Type 2)**
    - $(L_1 = \{x_1 \dots x_n \in \{0, 1\}^* \mid \sum_{k=1}^n x_{n-k+1} (\frac{3}{2})^{-k} < 1\} \in 1\text{ANNs} \setminus \text{CFLs})$
  - **1ANNs** with “**quasi-periodic**” weights  $\subseteq$  **FAs**  $\equiv$  **REG (Type 3)**
- the analog neuron hierarchy collapses at **3ANNs** (Šíma, 2018):
  - 3ANNs** = **4ANNs** = **5ANNs** =  $\dots \equiv$  **TM**s  $\equiv$  **RE (Type 0)**

# The Main Result: Separating 2ANNs From 1ANNs

“counting” language  $L_{\#} = \{0^n 1^n \mid n \geq 1\} \in \mathbf{2ANNs} \setminus \mathbf{1ANNs}$

$L_{\#}$  is a (non-regular) deterministic context-free language (DCFL)

accepted by a deterministic push-down automaton (DPDA)

## 1. $L_{\#} \notin \mathbf{1ANNs}$ :

**Theorem 1.** *The deterministic context-free language  $L_{\#}$  cannot be recognized by a neural network 1ANN with one extra analog unit having real weights.*

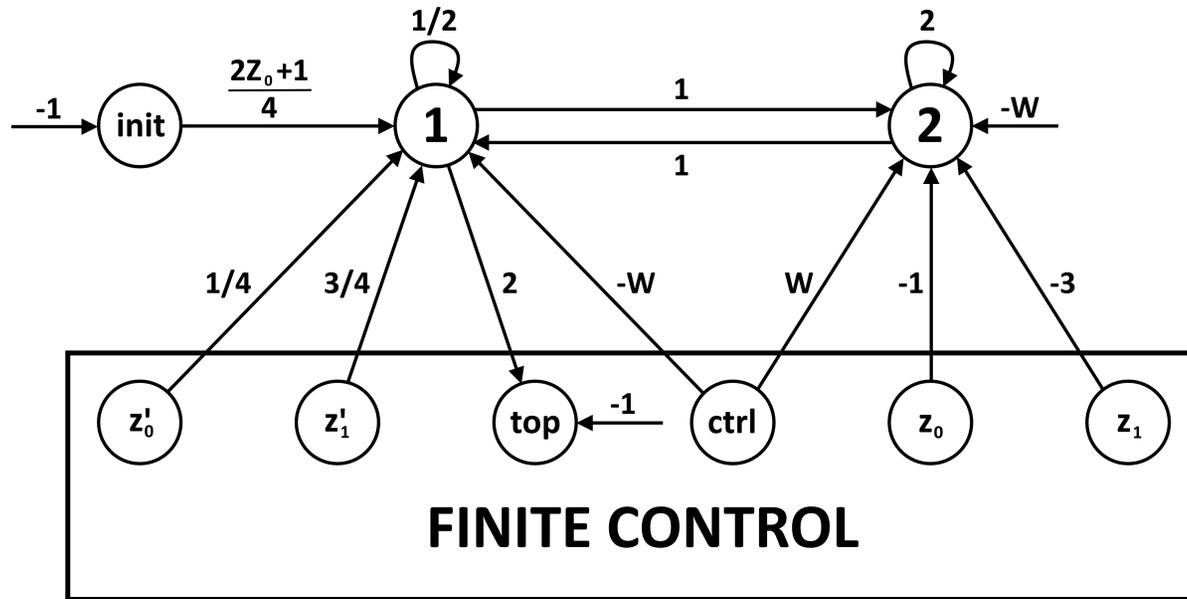
generalizes to  $(\mathbf{DCFLs} \setminus \mathbf{REG}) \cap \mathbf{1ANNs} = \emptyset$

i.e.  $\mathbf{1ANNs} \cap \mathbf{DCFLs} = \mathbf{0ANNs}$  (Šíma, Plátek, 2019)

## 2. $L_{\#} \in \mathbf{DCFLs} \subset \mathbf{2ANNs}$ :

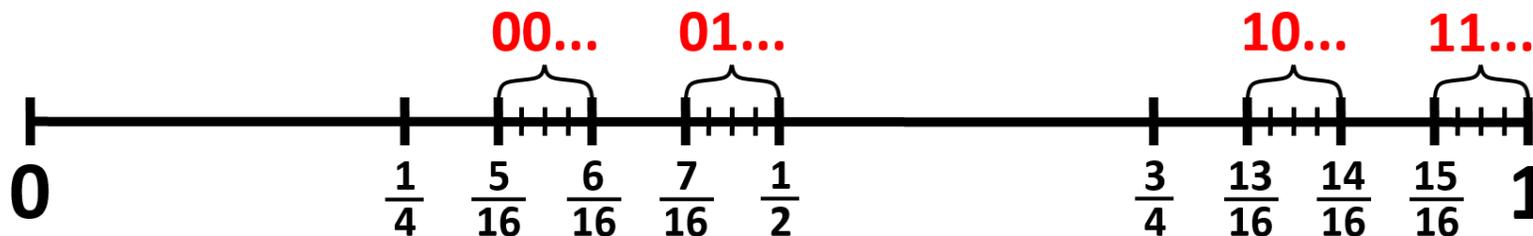
**Theorem 2.** *For any deterministic context-free language  $L \subseteq \{0, 1\}^*$ , there is a neural network 2ANN with two extra analog units having rational weights,  $\mathcal{N}$ , which accepts  $L = \mathcal{L}(\mathcal{N})$ .*

# A Schema of 2ANNs Simulating DPDAs



the stack contents  $x_1 \dots x_n \in \{0, 1\}^*$  are **encoded** by states  $y_1, y_2 \in [0, 1]$  of analog neurons 1 (push), 2 (pop) using **Cantor-like set** (Siegelmann, Sontag, 1995):

$$\text{code}(x_1 \dots x_n) = \sum_{i=1}^n \frac{2x_i + 1}{4^i} \in [0, 1]$$

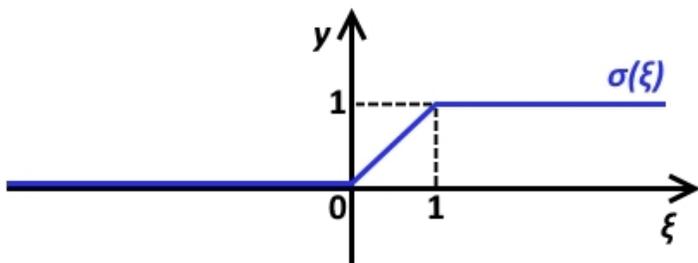


# Implementing the Stack Operations

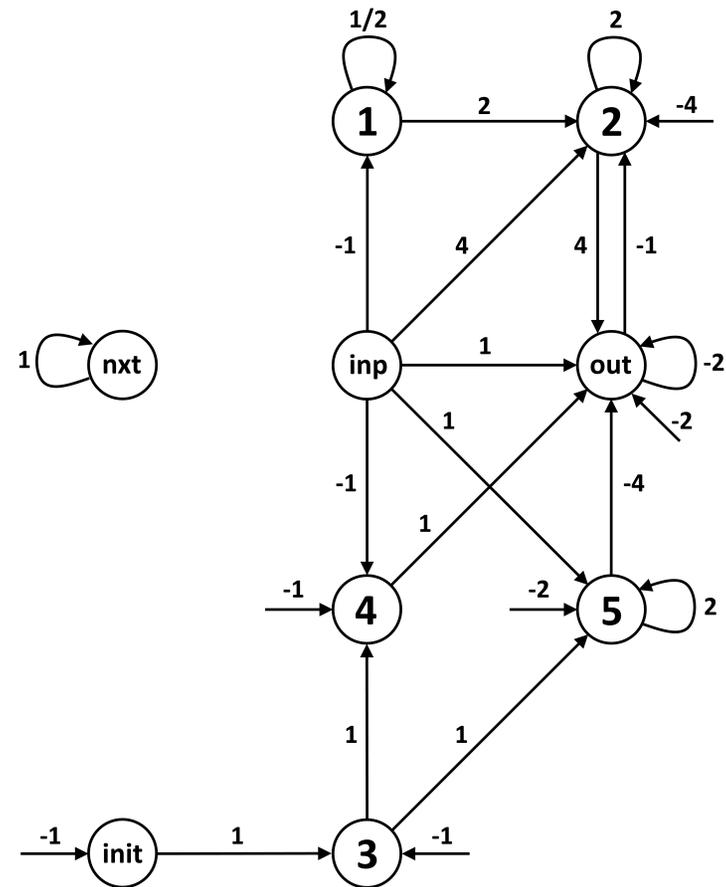
$$\text{top}(y_1) = H(2y_1 - 1) = \begin{cases} 1 & \text{if } y_1 \geq \frac{1}{2} \text{ i.e. } y_1 = \text{code}(1 x_2 \dots x_n) \\ 0 & \text{if } y_1 < \frac{1}{2} \text{ i.e. } y_1 = \text{code}(0 x_2 \dots x_n) \end{cases}$$

$$\begin{aligned} \text{pop}(y_2) &= \sigma(4y_2 - 2 \text{top} - 1) \\ &= \text{code}(x_2 \dots x_n) \end{aligned}$$

$$\begin{aligned} \text{push}(y_1, b) &= \sigma\left(\frac{1}{4}y_1 + \frac{2b+1}{4}\right) \\ &= \text{code}(b x_1 \dots x_n) \\ &\text{for } b \in \{0, 1\} \end{aligned}$$



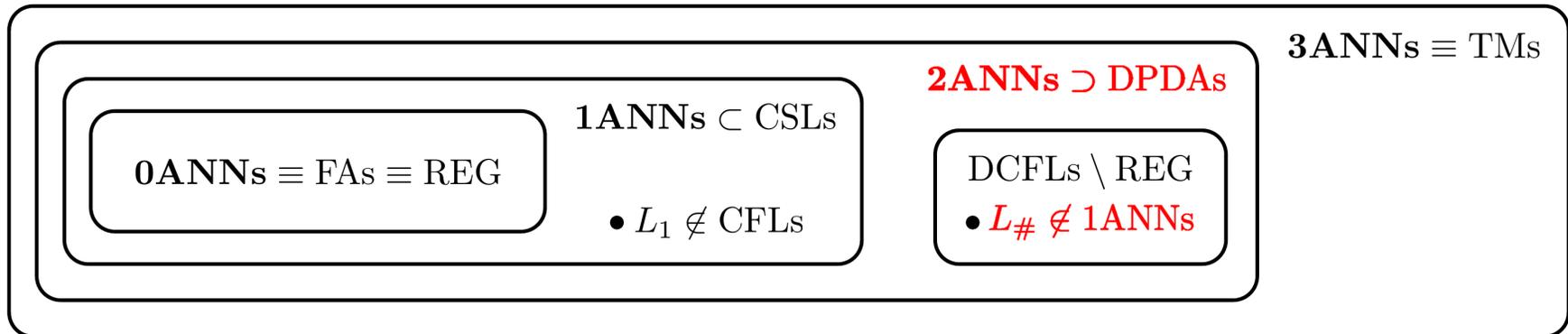
+ synchronizing the swaps  
between  $y_1$  and  $y_2$



Example of a 2ANN recognizing  $L_{\#}$

# A Summary of the Analog Neuron Hierarchy

$$\text{FAs} \equiv \text{0ANNs} \subsetneq \text{1ANNs} \subsetneq \text{2ANNs} \subseteq \text{3ANNs} \equiv \text{TM}s$$



## Open Problems:

- the separation of the 3rd level:  $2\text{ANNs} \subsetneq 3\text{ANNs} ?$
- strengthening the 2nd level separation to the **nondeterministic** CFLs:
 
$$(\text{CFLs} \setminus \text{REG}) \cap 1\text{ANNs} = \emptyset ?$$
- a proper “natural” hierarchy of NNs between integer and rational weights which can be mapped to known infinite hierarchies of REG/CFLs ?